

Exploring Hamiltonian Graphs: Theory and Applications

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ABSTRACT:

Hamiltonian graphs, a fundamental concept in graph theory, have attracted significant attention due to their rich structure and diverse applications in various fields such as computer science, network analysis, and optimization problems. This paper presents a comprehensive view of Hamiltonian graphs, focusing on their properties, characteristics for applications in real-world scenarios. We delve into the theoretical foundations of Hamiltonian graphs, discussing the necessary and sufficient conditions for a graph to be Hamiltonian, along with key theorems and concepts associated with Hamiltonian cycles and paths. Through this paper, we aim to provide a comprehensive understanding of Hamiltonian graphs and their significance in theoretical and practical contexts.

KEYWORDS:

Hamiltonian graph; Hamiltonian cycles; Hamiltonian paths; applications

1. Introduction

The study of graphs, which are mathematical structures made up of vertices (nodes) and edges (connections) connecting them, is the focus of the mathematical field of graph theory [1]. Hamiltonian graphs are one of the most important kinds of graphs because of their extensive structure and application. A graph that has a Hamiltonian cycle, or a cycle that visits every vertex precisely once aside from the starting and ending vertices, is said to be a Hamiltonian graph. Conversely, a path that makes one exact visit to each vertex is referred to as a Hamiltonian path. [2]. For a number of reasons, Hamiltonian graphs have attracted a lot related fields.

Hamiltonian graphs are tremendous theoretical interest. Their research contributes to our understanding of the underlying structure and characteristics of graphs in general, as they represent a class of graphs with particular traits [3]. Several key theorems and ideas in graph theory have been developed as a result of research on Hamiltonian graphs.

Hamiltonian graphs may be used to construct and solve a wide range of combinatorial optimisation problems [4]. For instance, the Travelling Salesman Problem (TSP), which has applications in

scheduling, logistics, and transportation, entails determining the shortest Hamiltonian cycle in a weighted graph. In the study of networks, identifying Hamiltonian cycles or paths can provide insights into the connectivity and traversal patterns within the network. This is crucial in fields such as telecommunications, computer networks, and social network analysis. Hamiltonian paths are relevant in the design and layout of integrated circuits, where the goal is to find a path that visits each component exactly once [5]. This aids in optimizing the wiring and layout of electronic components, leading to more efficient and compact circuit designs. Hamiltonian graphs have applications in game theory, particularly in the analysis of certain types of games and puzzles. Understanding the properties of Hamiltonian cycles and paths can provide insights into the strategies and optimal solutions for such games. The conditions for a graph to be Hamiltonian have been studied extensively, and several necessary and sufficient conditions have been established. Two well-known theorems that provide necessary and sufficient conditions for Hamiltonicity are Dirac's theorem and Ore's theorem.

1.1. Dirac's Theorem [2]

Dirac's theorem gives a necessary and sufficient condition for a graph to be Hamiltonian in terms of vertex degrees.

Let $(G = (V, E))$ be a simple graph with (n) vertices $(n \geq 3)$. If every vertex in (G) has degree at least $\frac{n}{2}$, then (G) is Hamiltonian.

Put differently, a graph (G) is required to be Hamiltonian if it is simple (i.e., does not contain loops or multiple edges) and each vertex has a degree that is at least half of the total number of vertices.

1.2. Ore's Theorem [3]

Another necessary and sufficient condition for Hamiltonicity, based on the sum of degrees of non-adjacent vertices, is provided by Ore's theorem.

Assume that the graph $(G = (V, E))$ has (n) vertices $(n \geq 3)$. (G) is Hamiltonian if the sum of the degrees of any pair of non-adjacent vertices (u) and (v) in (G) is larger than or equal to (n) .

In simple terms, if the sum of degrees of every pair of non-adjacent vertices in a graph is at least (n) , where (n) is the number of vertices, then the graph is guaranteed to be Hamiltonian.

Although Ore's and Dirac's theorems provide necessary and sufficient conditions for Hamiltonicity, it's crucial to keep in mind that their strict requirements may make practical application challenging. Though there are more requirements and conditions put forward by other scholars, the Dirac and Ore theorems are among the most well-known and often applied.

Our objective with regard to this work is to present a thorough grasp of Hamiltonian graphs and their importance in both theoretical and real-world settings.

2. Preliminaries

In graph theory, a unique kind of graph with a Hamiltonian cycle is called a Hamiltonian graph. A Hamiltonian graph has the following formal definition:

If a graph $(G = (V, E))$ has a Hamiltonian cycle, which is a cycle that visits every vertex precisely once except from the starting and ending vertices, which are the identical, then the graph is said to be Hamiltonian. Stated differently, a closed walk that begins at the beginning vertex and finishes precisely once at each vertex on a graph is known as a Hamiltonian cycle.

A Hamiltonian cycle in a graph (G) may be mathematically expressed as a permutation of the vertices in such a way that the permutation has an edge connecting each subsequent vertex and an edge connecting the initial and last vertices.

A graph is deemed Hamiltonian if there is such a permutation for it. Remember that not every graph has a Hamiltonian structure. A graph's ability to support a Hamiltonian cycle is dependent upon its particular connectivity and structure. One of the basic problems in graph theory is figuring out if a given graph is Hamiltonian [6]. Several criteria and techniques have been proposed to solve this problem. To summarize, a graph that comprises a Hamiltonian cycle, which visits every vertex precisely once, with the exception of the starting and ending vertices, which are identical, is referred to as a Hamiltonian graph. as shown in Fig 1

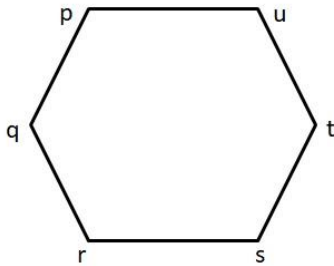


Figure 1. Hamiltonian graph

Definition 2.1:

A graph is considered planar if it can be plotted on a plane without any edge crossings. Put differently, the vertices can be represented as points and the edges as curves that link the points on a plane so that the only places where two edges overlap are at their ends [7].

3. Hamiltonian graph classes.

Hamiltonian graphs can belong to various classes based on their structural properties. Two important classes of Hamiltonian graphs are bipartite graphs and planar graphs.

3.1. Bipartite Graphs

3.1.1. Hamiltonicity in Bipartite

One unique trait of bipartite graphs is that they cannot have cycles of odd length. Because a Hamiltonian cycle must visit each vertex precisely once and have an even number of vertices (because it forms a cycle), a bipartite graph with more than two vertices is never Hamiltonian except for the complete bipartite graphs[8]. On the other hand, if a bipartite network contains just two vertices, it can be Hamiltonian.

Example 3.1:

If and only if ($m = n$), a complete bipartite graph ($K_{m,n}$), where (m) and (n) are the number of vertices in each partite set, is Hamiltonian.

Proof: Let's say that $m \neq n$. Let H represent the Hamiltonian that passes through each vertex in ($K_{m,n}$). $H = v_0e_0v_1e_1\dots v_ie_i$. The cycle needs to switch between the vertices on each side since K is bipartite. There is a $v_a = v_b, a < b$ within cycle H since $m \neq n$. Given that a cycle cannot have vertices that recur, this results in a contradiction. Therefore, if and only if $m = n$, a complete bipartite graph $K_{m,n}$ contains a Hamilton cycle.

3.1.2. Hamiltonicity in Planar Graphs

Determining the Hamiltonian nature of a planar graph is usually a difficult job. Planar graphs are not always Hamiltonian, even though they can be. The topic of determining the necessary and sufficient criteria for Hamiltonicity in planar graphs is still unresolved as there exist non-Hamiltonian planar graphs[9].

Example 3.2: The graph of a cube (hexahedron) is a planar graph that is Hamiltonian.

In graph theory, bipartite and planar graphs are significant graph types. To fully explore the field of Hamiltonian graphs and their applications, one must have a solid knowledge of their Hamiltonicity features.

4. Applications of Hamiltonian Graphs

4.1. Network routing: Hamiltonian cycles in communication networks.

Hamiltonian cycles have a major impact on network routing, especially when it comes to communication networks, since they guarantee dependable and effective data transfer. The relevance of Hamiltonian cycles in network routing is as follows:

4.1.1. Connectivity and Redundancy

Hamiltonian cycles provide a loop that goes around each node precisely once, ensuring connectedness inside the network. No matter how big or complicated the network is, data must be able to flow between any pair of nodes for there to be connection. Additionally, since data may still be sent through alternate channels in the event of a connection failure, the redundancy offered by Hamiltonian cycles improves fault tolerance.

4.1.2. Routing Protocol Optimization

Communication networks can benefit from the use of Hamiltonian cycles to optimise their routing methods. Routing algorithms may be engineered to effectively distribute traffic around the network, minimising delays and congestion, by guaranteeing that every node is a member of a Hamiltonian cycle. Furthermore, packets can be routed along preset courses specified by Hamiltonian cycles, which streamlines the routing process.

4.1.3. Multicast and Broadcast Routing

Hamiltonian cycles are particularly useful in multicast and broadcast routing scenarios, where data needs to be transmitted to multiple destinations simultaneously or to all nodes in the network. By utilizing Hamiltonian cycles, multicast trees can be constructed to efficiently deliver data to multiple recipients, ensuring that each node receives the message exactly once without unnecessary duplication or loops.

4.1.4. Load Balancing and Traffic Engineering

In communication networks, load balancing and traffic engineering may be made easier with the use of Hamiltonian cycles. Network resources may be used more effectively by spreading traffic equally over the Hamiltonian cycle's edges, avoiding bottlenecks and congestion at particular nodes or links. Furthermore, Hamiltonian cycles offer an organized framework for network-wide bandwidth and resource allocation optimization [10].

4.1.5. Scalability and Expansion

Hamiltonian cycles provide a scalable framework for network expansion and growth. As the network evolves and new nodes are added, Hamiltonian cycles can be dynamically adjusted or extended to accommodate the changes, ensuring continued connectivity and efficiency. This scalability is essential for accommodating the increasing demands and complexities of modern communication networks.

Overall, Hamiltonian cycles play a crucial role in network routing by providing connectivity, redundancy, and optimization mechanisms that ensure efficient and reliable data transmission in communication networks. By leveraging the properties of Hamiltonian cycles, network designers and engineers can develop robust routing solutions that meet the evolving needs of today's interconnected world.

4.2. Circuit design: Hamiltonian paths in VLSI design.

Hamiltonian routes are important in many aspects of circuit design in VLSI (Very Large Scale Integration) design, especially in the layout and routing stages. The following explains why Hamiltonian routes matter in VLSI design:

4.2.1. Routing Optimization

Routing in VLSI entails figuring out the pathways that connect the chip's components (logic gates, memory cells, etc.) during the physical design stage. Hamiltonian routes offer a path that visits each component precisely once, which may be used to optimise routing. Designers can minimise power consumption and signal propagation delays by minimising the length of the routing cables through the discovery of Hamiltonian pathways.

4.2.2. Layout Compactness

By guaranteeing that each component can be accessed from every other component via a single path, Hamiltonian pathways help to make the chip architecture compact. This results in denser and more effective architectures by lowering the total area taken up by routing cables and interconnects. Hamiltonian routes also aid in the organised organisation of the layout, streamlining the design process and enhancing manufacturability.

4.2.3. Avoiding Signal Crosstalk

Signal crosstalk, which happens when signals from nearby lines interfere with one another, may be prevented with the use of Hamiltonian routes. In order to reduce the possibility of crosstalk and signal deterioration, designers can make sure that wires carrying similar signals are routed separate from one another by adhering to Hamiltonian routes during routing.

4.2.4. Testing and Debugging

Hamiltonian paths can facilitate testing and debugging of VLSI circuits by providing a systematic way to traverse the components on the chip. During testing, test patterns can be applied sequentially along the Hamiltonian path, ensuring comprehensive coverage of the circuit. Similarly, in the debugging phase, designers can trace the behavior of signals along the Hamiltonian path to identify and diagnose potential issues or faults.

4.2.5. Scalability and Flexibility

Hamiltonian paths provide a scalable and flexible framework for VLSI design, allowing for efficient expansion and modification of the circuit layout. As the complexity of VLSI designs continues to increase, Hamiltonian paths offer a structured approach to routing that can accommodate the growing number of components and interconnects while maintaining performance and reliability.

Overall, Hamiltonian paths play a crucial role in VLSI design by optimizing routing, improving layout compactness, reducing signal crosstalk, facilitating testing and debugging, and ensuring scalability and flexibility. By leveraging the properties of Hamiltonian paths, VLSI designers can develop efficient and robust integrated circuits that meet the demanding requirements of modern electronic systems.

4.3. Other real-world applications.

4.3.1. Robotics & Path Planning

Hamiltonian pathways are useful in robotics to find the best routes for robots to travel across surroundings while stopping at each spot precisely once. This is especially true when it comes to autonomous navigation and path planning. Applications like driverless cars, robotic exploration missions, and warehouse automation depend on this application.

4.3.2. Wireless Sensor Networks

Hamiltonian cycles are used to collect and route data in wireless sensor networks in an effective manner. Data may be reliably gathered from each sensor node and communicated to a central sink node by arranging the nodes into a Hamiltonian cycle. This minimises energy usage and increases the lifetime of the network.

4.3.3. Genetic techniques and Optimization

The creation of optimization techniques, including simulated annealing and genetic algorithms, is inspired by Hamiltonian cycles and routes. In order to discover the best answers to challenging optimization issues in a variety of fields, such as network design, scheduling, and resource allocation, these algorithms imitate the process of evolution.

4.3.4. Circuit Testing and Fault Detection

To guarantee thorough coverage of the circuit under test, Hamiltonian cycles are used in circuit testing and fault detection processes. By identifying errors and confirming integrated circuit performance, test patterns created using Hamiltonian cycles contribute to the dependability and high calibre of electronic equipment.

4.4. Hamiltonian graph's application in chemistry

Chemistry uses Hamiltonian graphs for studying chemical structures and reactions. In general Hamiltonian graphs have a number of uses:

4.4.1. Chemical Graph Theory

Chemical compounds may be visualised as graphs, with bonds serving as the edges and atoms as the vertices. These graphs' Hamiltonian routes or cycles can provide information about the connection and structure of molecules.

4.4.2. Reaction Pathways

Knowing the potential routes by which reactants might change into products is essential for analyzing chemical processes. Reaction graphs using Hamiltonian cycles can show possible reaction routes.

4.4.3. Organic Chemistry

Organic molecule reactivity and structure may be investigated using Hamiltonian cycles. For instance, Hamiltonian cycles can inform the research of cyclic molecules such as cycloalkanes and aromatic compounds about potential isomers and their stability.

4.4.4. Topology and Connectivity

Molecular structural topology and connectivity can be elucidated using Hamiltonian graphs. Understanding the stability and reactivity of complicated compounds is made easier by this.

4.4.5. Computational Chemistry

Graph-based algorithms are widely used in computational chemistry to model and simulate molecules. Reaction pathway prediction and chemical structure optimisation are two applications of Hamiltonian graphs.

4.4.6. Chemical Graph Isomorphism

The problem of chemical graph isomorphism, which asks if two molecular graphs reflect the same chemical molecule, may also be solved using the study of Hamiltonian graphs.

5. Emerging applications and research directions

Combining uses and avenues for future study in Hamiltonian graph theory span many different academic domains and technological developments. As quantum computing gains popularity, researchers are becoming more interested in studying Hamiltonian graphs and their characteristics in

relation to quantum information processing and quantum algorithms. A new area of study is looking at quantum analogues of classical Hamiltonian graph problems and creating quantum algorithms for effective solution.

Hamiltonian graphs can help with power distribution optimization, transmission loss minimization, and grid resilience enhancement in the context of smart grids and energy networks. The advancement of smart grid technology depends on research into the modeling and analysis of energy networks as Hamiltonian graphs as well as the creation of effective routing and scheduling algorithms.

Protein interactions, genetic processes, and molecular structures may all be modeled using Hamiltonian graphs in these fields. Graph-based metabolic pathway analysis, genome sequencing, and the investigation of Hamiltonian cycles in biological networks are examples of new research topics with potential applications in personalized medicine and medication development.

Using Hamiltonian graph theory to analyse social networks and influence dynamics can provide light on the spread of information, the formation of opinions, and the identification of communities. Computational social science benefits from research on influence propagation mechanisms, viral spread prediction, and social network modeling as Hamiltonian graphs. The design and optimization of these networks both benefit from an understanding of Hamiltonian graph theory. Emerging research approaches having implications for IoT applications in smart cities, healthcare, and environmental monitoring include examining Hamiltonian routes for effective data collecting, routing protocols for IoT devices, and network optimization strategies for energy-efficient communication.

Hamiltonian graphs provide information on task distribution, path planning techniques, and swarm coordination in the context of robotic swarm systems. An emerging field with applications in autonomous robotics and swarm intelligence is the construction of decentralized algorithms for building Hamiltonian cycles in robotic swarms, optimizing collective behaviors, and accomplishing global objectives through local interactions. In addition to advancing the theoretical underpinnings of graph theory, investigating these new applications and research avenues in Hamiltonian graph theory tackles practical issues in a variety of domains, spurring creativity and technological advancement. In order to fully use Hamiltonian graphs in solving intricate real-world issues, interdisciplinary cooperation and the integration of theoretical ideas with domain-specific expertise are crucial.

6. Conclusion

A key idea in graph theory, Hamiltonian graphs have garnered a lot of interest because of its varied applications in computer science, network analysis, and optimisation issues, as well as their complex structure. This work offers a thorough understanding of Hamiltonian graphs with an emphasis on their attributes and uses in practical contexts. We go into the theoretical underpinnings of Hamiltonian graphs, covering the essential theorems and ideas related to Hamiltonian cycles and routes as well as the necessary and sufficient criteria for a graph to be Hamiltonian. Our goal in writing this work is to present a thorough grasp of Hamiltonian graphs and their importance in both theoretical and real-world settings.

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